

# Nonconcave Utility Maximization

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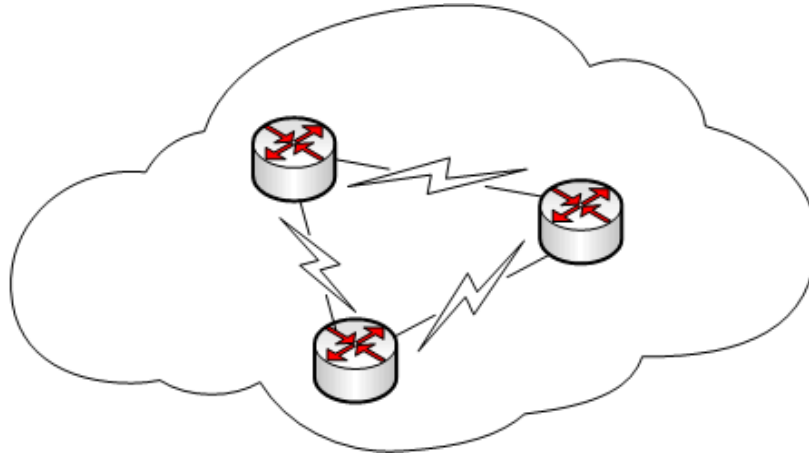


# Objectives

- Optimization problems motivated by practical applications
- Develop algorithms that can be “easily implemented”
- Decentralization is a must



# Motivation



Problem statement:

$$\max \sum_{i \in \mathcal{S}} U_i(r_i) \text{ s.t. network resource constraints}$$



# Motivation (cont.)

## Challenges:

1. IP video traffic will be 82% of all consumer Internet traffic by 2020, up from 70% in 2015<sup>1</sup>.
2. Large-scale networks.

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<sup>1</sup> ***Cisco VNI Forecast and Methodology, 2015-2020.***

## Assumptions:

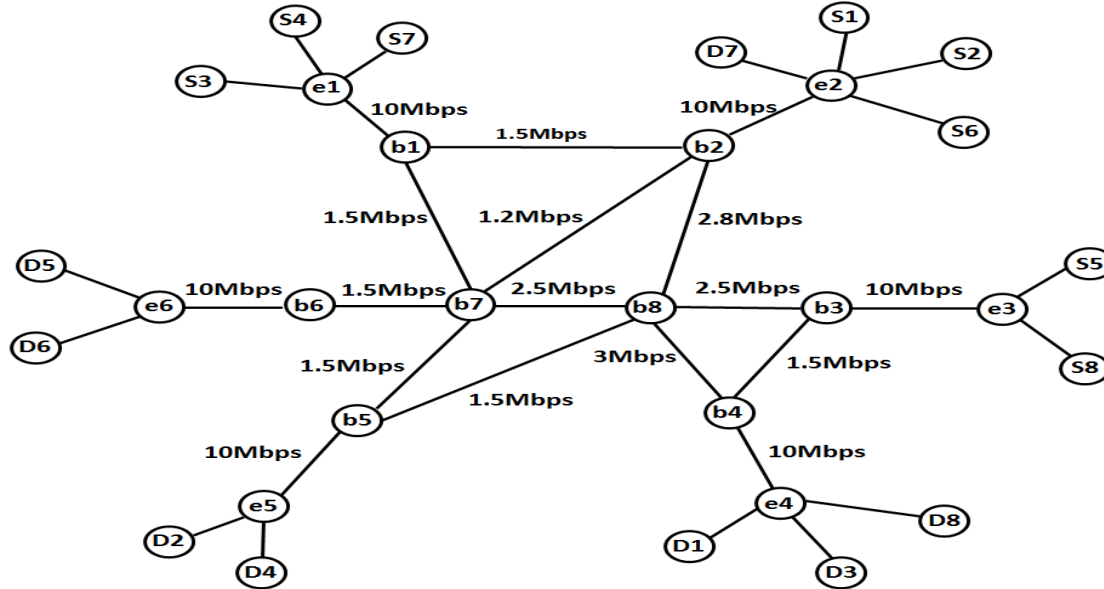
1. Nonconcave network utility functions<sup>2</sup>.
2. No global information available at any network entity.

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<sup>2</sup> ***Yin, X et al, A control-Theoretic Approach for Dynamic Adaptive Video Streaming over http, SIGCOMM, 2015.***



# Flow Control



## Notation:

$(\mathcal{S}, \mathcal{L})$  : set of sources and links, respectively.

$x_i^p$  : data rate of source  $i$  on path  $p$ ,  $p \subset \mathcal{L}$ .

$$\mathbf{x}_i = [x_i^p]_{p \in \mathcal{P}_i}, \quad r_i = \sum_{p \in \mathcal{P}_i} x_i^p.$$

$$\mathbf{A}_i = [a_{l,p}^i]_{l \in \mathcal{L}, p \in \mathcal{P}_i}.$$



# Problem Formulation

$$\begin{aligned} \max_{\mathbf{x}, \mathbf{r}} \quad & \sum_{i \in \mathcal{S}} U_i(r_i) \\ \text{s.t.} \quad & \sum_{i \in \mathcal{S}} \mathbf{A}_i \mathbf{x}_i \preceq \mathbf{c} \\ & r_i = \sum_{p \in \mathcal{P}_i} x_i^p \\ & 0 \leq b_i \leq r_i \leq B_i \\ & x_i^p \geq 0 \end{aligned}$$

Structure to be exploited:

- Utility function “decentralized”
- $a_{i,j} \geq 0$



# Class of Utility Functions

We consider utility functions of the form

$$U_i(r_i) = \sum_{j=0}^{\ell} p_{i,j} \left( r_i^{1/\ell} \right)^j$$

Motivation:

- “General set” of utility functions
- Leads to a formulation solved efficiently by distributed algorithms.



# Measures and Optimization

**Theorem.** *Let  $f$  be a real-valued function,  $\mathcal{G}$  be a compact set, and  $\mu$  be a probability measure with support  $\text{supp}(\mu)$ . Then,*

$$\inf_x \{f(x) : x \in \mathcal{G}\} = \inf_{\mu} \left\{ \int f d\mu : \text{supp}(\mu) \subset \mathcal{G} \right\}$$



# Moments of Univariate Measures

**Theorem.** Given  $\{m_j\}_{j=1}^\ell$ , there exists a Borel measure  $\mu(\cdot)$  with support contained in  $\mathcal{I} = [-\epsilon, \epsilon]$  such that  $\mu(\mathcal{I}) = 1$  and  $m_j = \int_{\mathcal{I}} y^j d\mu$  if and only if  $\mathbf{M}(0, \ell) \succeq 0$ ,  $\epsilon^2 \mathbf{M}(1, \ell - 1) \succeq \mathbf{M}(2, \ell)$ .

$$\mathbf{M}(k, k + 2h) = \begin{bmatrix} m_k & m_{k+1} & \dots & m_{k+h} \\ m_{k+1} & \ddots & \ddots & m_{k+h+1} \\ \vdots & \ddots & \ddots & \vdots \\ m_{k+h} & \dots & \dots & m_{k+2h} \end{bmatrix}$$



# Back to “our” Problem

$$\begin{aligned} \max_{\mathbf{x}, \mathbf{r}} \quad & \sum_{i \in \mathcal{S}} \sum_{j=0}^{\ell} p_{i,j} r_i^{j/\ell} \\ \text{s.t.} \quad & \sum_{i \in \mathcal{S}} \mathbf{A}_i \mathbf{x}_i \preceq \mathbf{c} \\ & r_i = \sum_{p \in \mathcal{P}_i} x_i^p \\ & 0 \leq b_i \leq r_i \leq B_i \\ & x_i^p \geq 0 \end{aligned}$$

$\Leftrightarrow$

$$\begin{aligned} \max_{\mathbf{x}, \mathbf{r}} \quad & \sum_{i \in \mathcal{S}} \sum_{j=0}^{\ell} p_{i,j} y_i^j \\ \text{s.t.} \quad & \sum_{i \in \mathcal{S}} \mathbf{A}_i \mathbf{x}_i \preceq \mathbf{c} \\ & 0 \leq y_i \leq r_i^{1/\ell} \\ & r_i = \sum_{p \in \mathcal{P}_i} x_i^p \\ & 0 \leq b_i \leq r_i \leq B_i \\ & x_i^p \geq 0 \end{aligned}$$



# Moment Formulation

$$\max_{\mathbf{m}, \mathbf{x}, \mathbf{r}} \sum_{i \in \mathcal{S}} \mathbf{p}_i^\top \mathbf{m}_i$$

$$\text{s.t. } \mathbf{x} \in \mathcal{C}$$

$$(\mathbf{m}_i, \mathbf{x}_i, r_i) \in \mathcal{K}_i, \quad i \in \mathcal{S}$$

$$\mathcal{C} = \left\{ \mathbf{x} : \sum_{i \in \mathcal{S}} \mathbf{A}_i \mathbf{x}_i \preceq \mathbf{c} \right\}$$

$$\mathcal{X}_i = \left\{ (\mathbf{x}_i, r_i) \in \mathbb{R}_+^{|\mathcal{P}_i|} \times \mathbb{R}_+ : r_i = \sum_{p \in \mathcal{P}_i} x_i^p, \quad b_i \leq r_i \leq B_i \right\}$$

$$\mathcal{K}_i = \left\{ (\mathbf{m}_i, \mathbf{x}_i, r_i) : \mathbf{M}_i(0, \alpha) \succeq 0, \quad B_i^2 \mathbf{M}_i(1, \alpha - 1) \succeq \mathbf{M}_i(2, \alpha), \right. \\ \left. m_{i,j} \leq r_i^{j/\ell}, \quad j = 1, \dots, \alpha, (\mathbf{x}_i, r_i) \in \mathcal{X}_i \right\}$$



# Comments on Moment Formulation

- Converges to the desired optimum as  $\alpha \rightarrow \infty$
- “Nice” separation between local and global constraints
- Convex if  $\alpha \leq \ell$



# Decentralized Implementation

**Algorithm 1:** Distributed Traffic Allocation Algorithm (DTTA)

$(\rho, \{\alpha^n\}_{n \in \mathbb{Z}_+}, \{\tau^k\}_{k \in \mathbb{N}}, \{\lambda^k\}_{k \in \mathbb{N}}, \mathbf{z}^0 = [\mathbf{z}_i^0]_{i \in \mathcal{S}}, \mathbf{u}^0 = [\mathbf{u}_i^0]_{i \in \mathcal{S}})$

Initialize  $\mathbf{z}^0, \mathbf{u}^0$

**for**  $k = 0, 1, \dots$  **do**

$(\mathbf{m}_i, \mathbf{x}_i, r_i)^{k+1} \leftarrow \operatorname{argmax} \left\{ \mathbf{p}_i^\top \mathbf{m}_i - \frac{\rho}{2} \|\mathbf{x}_i - \mathbf{z}_i^k + \mathbf{u}_i^k\|^2 : (\mathbf{m}_i, \mathbf{x}_i, r_i) \in \mathcal{K}_i \right\}, i \in \mathcal{S}.$

Initialize  $\mathbf{z}^{k,1} \leftarrow \mathbf{z}^k$

**for**  $n = 1, \dots, \tau^k - 1$  **do**

Each source  $i \in \mathcal{S}$  transmits data at rate  $(z_i^p)^{k,n}$  along the path  $p \in \mathcal{P}_i$ .

For each link  $l \in \mathcal{L}$ :  $b_l^{k,n} \leftarrow 1$  if link  $l$  is congested;  $b_l^{k,n} \leftarrow 0$  otherwise.

Each link  $l \in \mathcal{L}$  sends one bit  $b_l^{k,n}$  to all  $i \in \mathcal{S}_l$ .

$\mathbf{g}_i^{k,n} = \mathbf{z}_i^{k,n} - \mathbf{x}_i^{k+1} - \mathbf{u}_i^k + \lambda^k \sum_{l \in \mathcal{L}} b_l^{k,n} \mathbf{a}_l^i, i \in \mathcal{S}.$

$\mathbf{z}_i^{k,n+1} = \left( \mathbf{z}_i^{k,n} - \alpha^n \mathbf{g}_i^{k,n} \right)^+, i \in \mathcal{S}.$

$\mathbf{z}^{k+1} \leftarrow \mathbf{z}^{k,\tau^k}$

$\mathbf{u}_i^{k+1} \leftarrow \mathbf{u}_i^k + \mathbf{x}_i^{k+1} - \mathbf{z}_i^{k+1}, i \in \mathcal{S}.$



# Why does it work ?

Equivalent formulations:

$$\max_{\mathbf{m}, \mathbf{x}, \mathbf{r}} \left\{ \sum_{i \in \mathcal{S}} \mathbf{p}_i^\top \mathbf{m}_i : \mathbf{x} \in \mathcal{C}, (\mathbf{m}_i, \mathbf{x}_i, r_i) \in \mathcal{K}_i, i \in \mathcal{S} \right\}$$

$$\max_{\mathbf{m}, \mathbf{x}, \mathbf{r}, \mathbf{z}} \left\{ \sum_{i \in \mathcal{S}} \mathbf{p}_i^\top \mathbf{m}_i : \mathbf{z} \in \mathcal{C}, (\mathbf{m}_i, \mathbf{x}_i, r_i) \in \mathcal{K}_i, \mathbf{z}_i = \mathbf{x}_i, i \in \mathcal{S} \right\}$$

Augmented Lagrangian:

$$L(\mathbf{m}, \mathbf{x}, \mathbf{r}, \mathbf{z}, \nu) = \sum_{i \in \mathcal{S}} [\mathbf{p}_i^\top \mathbf{m}_i - \nu_i^\top (\mathbf{x}_i - \mathbf{z}_i) - (\rho/2) \|\mathbf{x}_i - \mathbf{z}_i\|^2]$$

$$g(\nu) = \max_{\mathbf{m}, \mathbf{x}, \mathbf{r}, \mathbf{z}} \{L(\mathbf{m}, \mathbf{x}, \mathbf{r}, \mathbf{z}, \nu) : (\mathbf{m}_i, \mathbf{x}_i, r_i) \in \mathcal{K}_i, i \in \mathcal{S}, \mathbf{z} \in \mathcal{C}\}$$



# Why does it work? (cont.)

ADMM:

$$(\mathbf{m}, \mathbf{x}, \mathbf{r})^{k+1} = \operatorname{argmax}_{\mathbf{m}, \mathbf{x}, \mathbf{r}} \left\{ L(\mathbf{m}, \mathbf{x}, \mathbf{r}, \mathbf{z}^k, \nu^k) : (\mathbf{m}_i, \mathbf{x}_i, r_i) \in \mathcal{K}_i, i \in \mathcal{S} \right\}$$

$$\mathbf{z}^{k+1} = \operatorname{argmax}_{\mathbf{z}} \left\{ L(\mathbf{m}^{k+1}, \mathbf{x}^{k+1}, \mathbf{r}^{k+1}, \mathbf{z}, \nu^k) : \mathbf{z} \in \mathcal{C} \right\}$$

$$\nu^{k+1} = \nu^k + \rho(\mathbf{x}^{k+1} - \mathbf{z}^{k+1})$$

Decentralization:

$$(\mathbf{m}_i, \mathbf{x}_i, r_i)^{k+1} = \operatorname{argmax}_{(\mathbf{m}_i, \mathbf{x}_i, r_i) \in \mathcal{K}_i} \left\{ \mathbf{p}_i^\top \mathbf{m}_i - (\rho/2) \|\mathbf{x}_i - \mathbf{z}_i^k + \mathbf{u}_i^k\|^2 \right\}$$

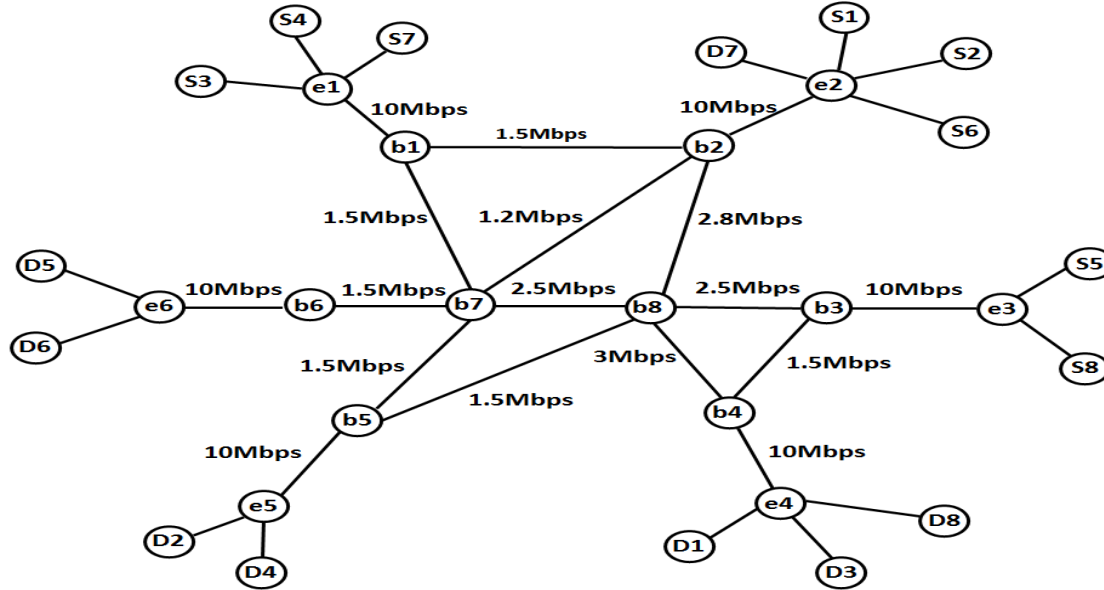
$$\mathbf{z}^{k+1} = \Pi_{\mathcal{C}} \left\{ \mathbf{x}^{k+1} + \mathbf{u}^k \right\}$$

$$\mathbf{u}_i^{k+1} = \mathbf{u}_i^k + \mathbf{x}_i^{k+1} - \mathbf{z}_i^{k+1}$$





# Simulations



Parameter values:

$$\ell = 6$$

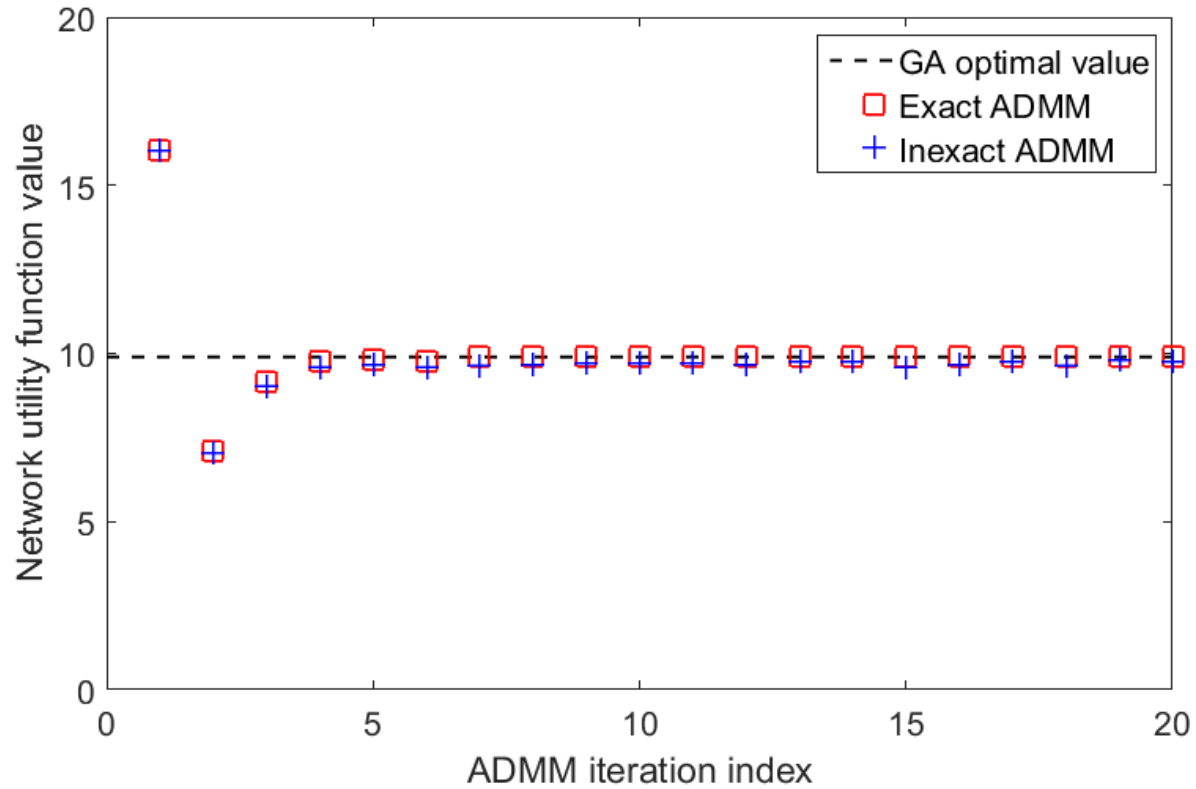
$$b_i = 0, B_i = 10$$

$$\rho = 1$$

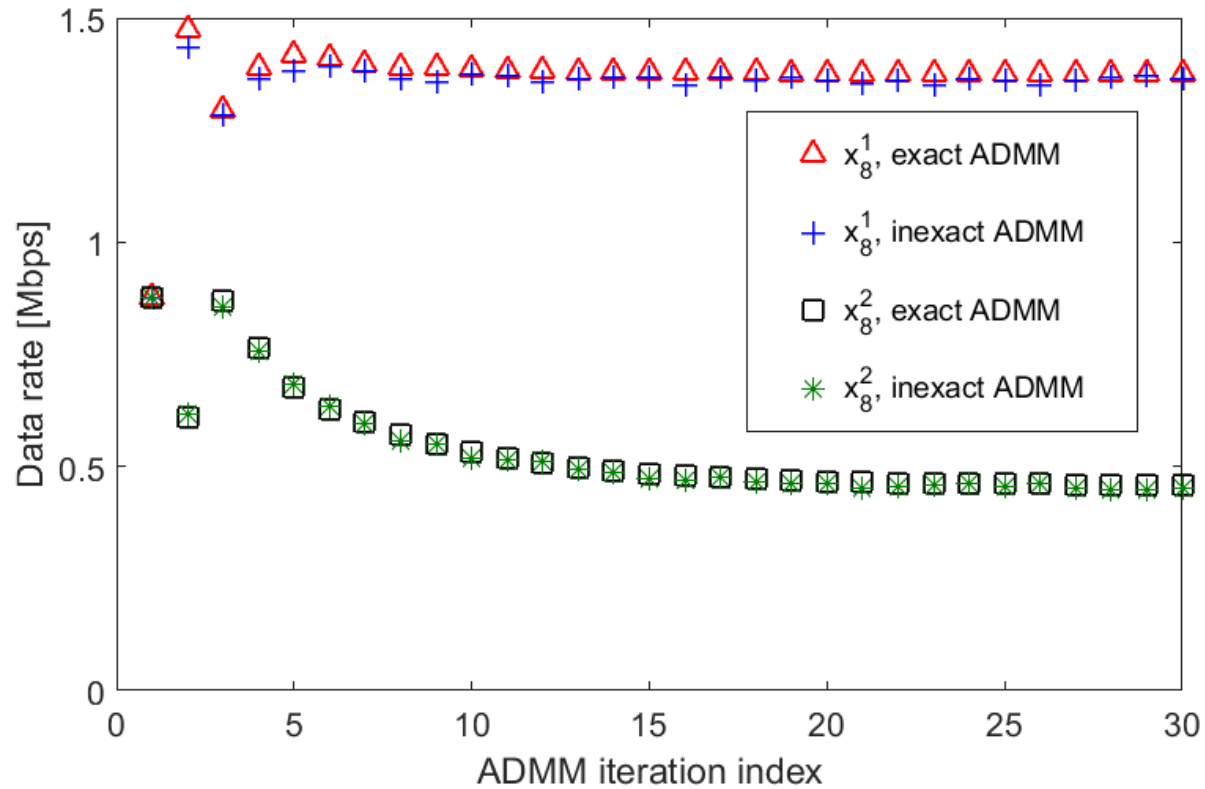
$$\lambda^k = 10 \forall k \in \mathbb{N}, \tau^k = 10^3 \forall k \in \mathbb{N}, \alpha^n = 1/n$$



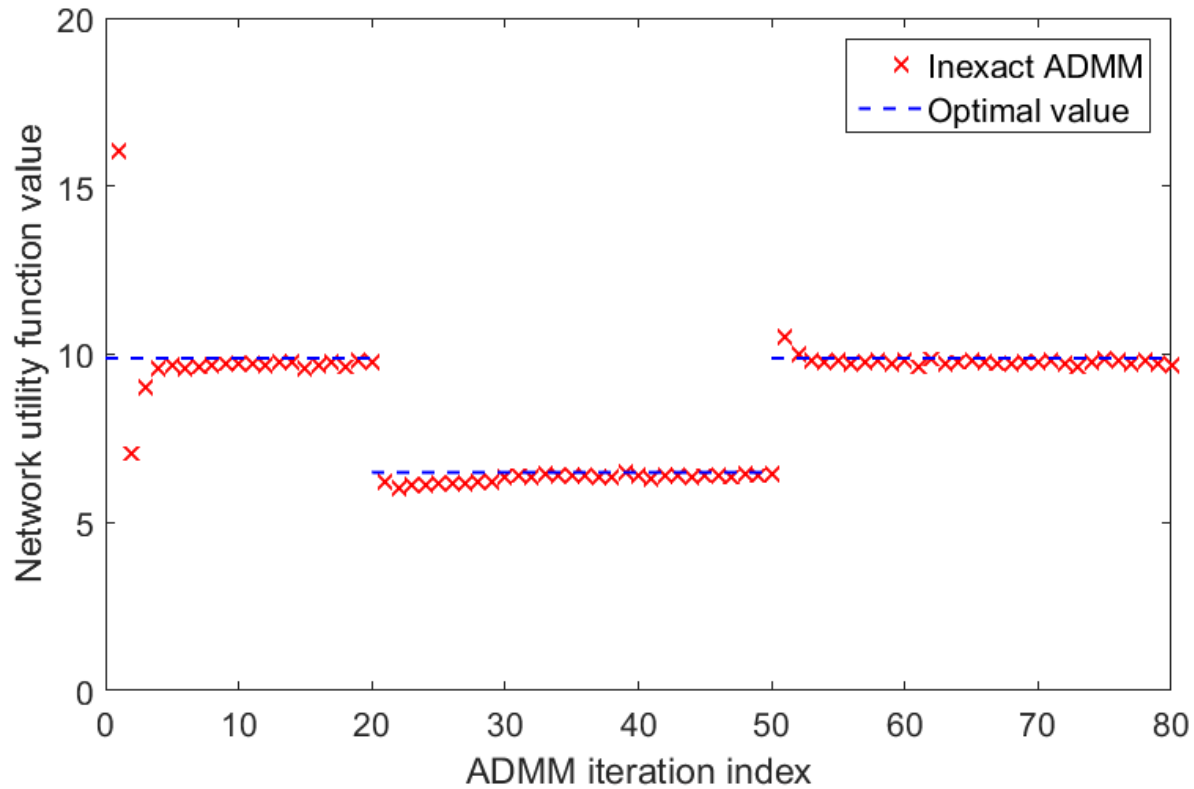
# Network Utility



# Data Rate Allocation



# Robustness to Link Failures



# Concluding Remarks

- General approach to utility maximization
- Efficiently solution of a class of polynomial optimization problems
- Proposed algorithm is decentralized and can solve large problems

Further work:

- What other classes of polynomials optimization problems can be efficiently solved?



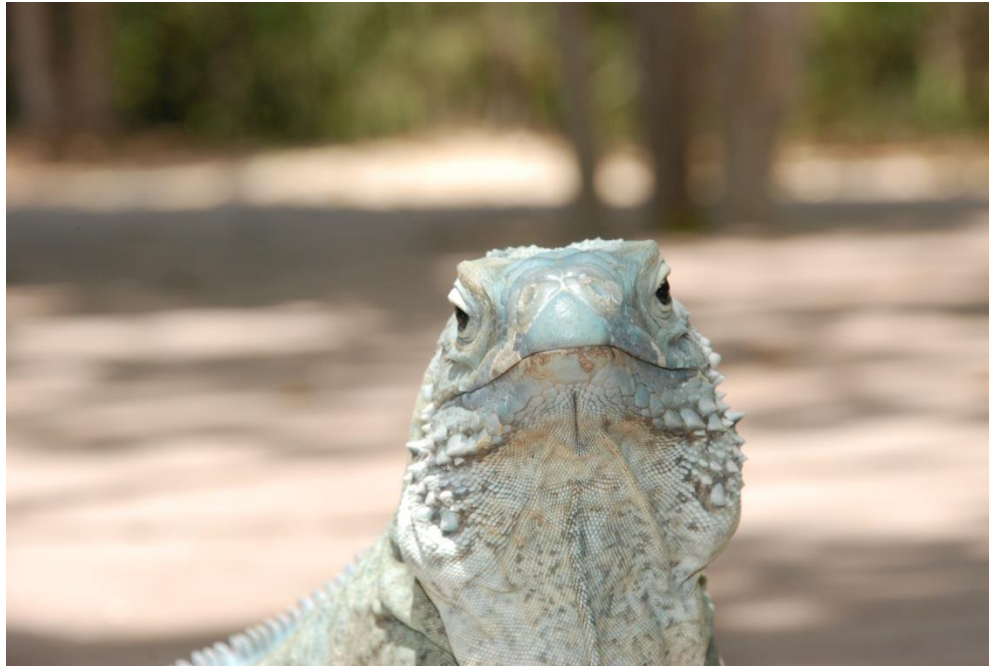
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Thank you

Merci beaucoup



Questions?

